



MiniZinc

Technologie programistyczne, 2018-2019

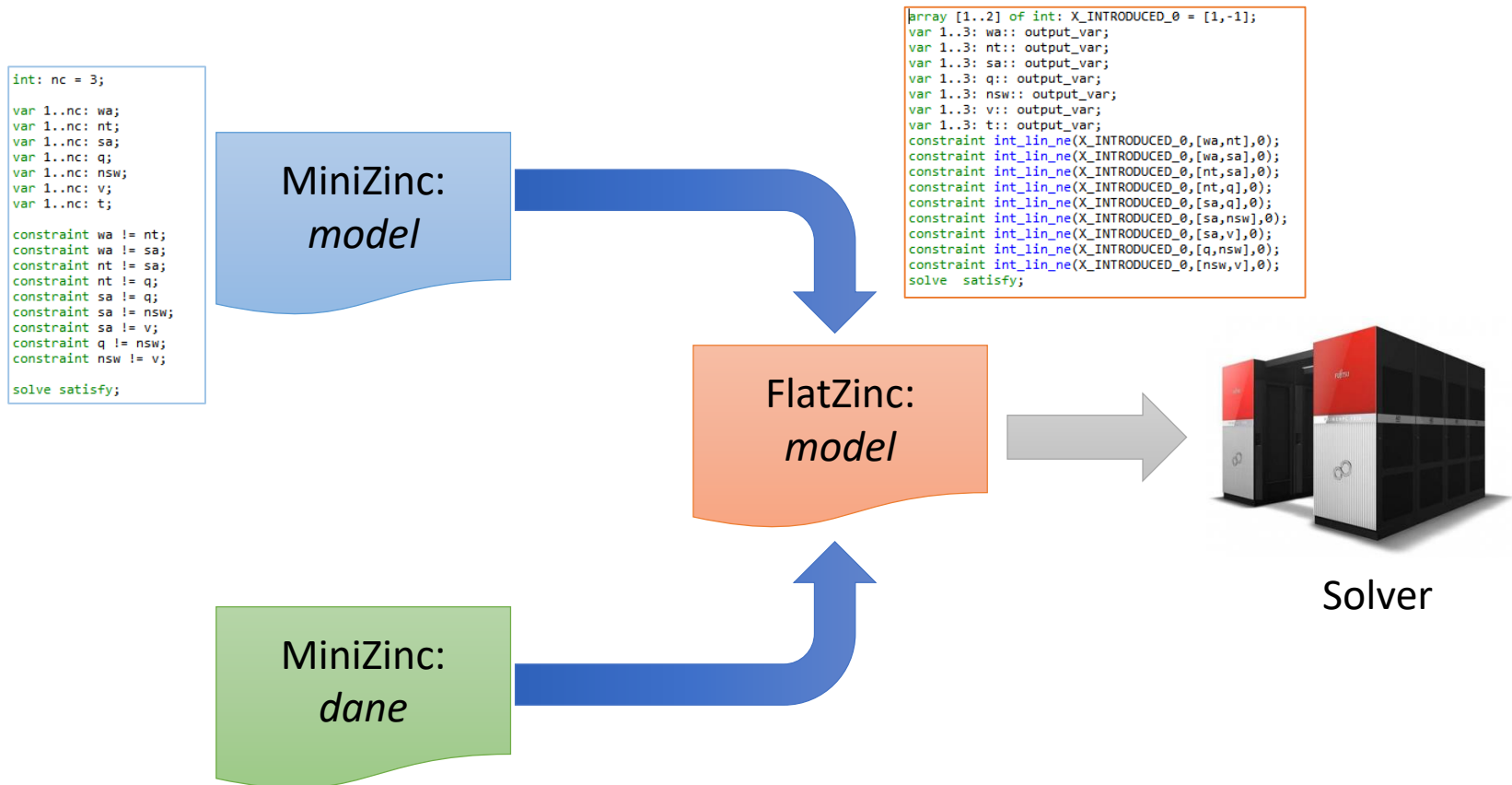
<http://www.minizinc.org>

Co to jest MiniZinc?

- Język do opisów problemów decyzyjnych i optymalizacyjnych z ograniczeniami
- Zmienne decyzyjne liczbowe (`int`, `float`), logiczne (`bool`) oraz wyliczeniowe (`enum`)
- Możliwość współpracy z wieloma specjalizowanymi solverami, także implementacje referencyjne
- Rozwijany przez Monash University, Data61 Data Science i University of Melbourne (Australia)



MiniZinc i FlatZinc



Problem 1: kolorowanie mapy



```
1 int: nc = 3;
```

```
3 var 1..nc: wa;  
4 var 1..nc: nt;  
5 var 1..nc: sa;  
6 var 1..nc: q;  
7 var 1..nc: nsw;  
8 var 1..nc: v;  
9 var 1..nc: t;
```

```
11 constraint wa != nt;  
12 constraint wa != sa;  
13 constraint nt != sa;  
14 constraint nt != q;  
15 constraint sa != q;  
16 constraint sa != nsw;  
17 constraint sa != v;  
18 constraint q != nsw;  
19 constraint nsw != v;
```

```
21 solve satisfy;
```

```
23 output ["wa=\(wa) nt=\(nt) sa=\(sa) q=\(q) nsw=\(nsw) v=\(v) t=\(t)"];
```

parametr – jednokrotne przypisanie wartości
typy parametrów: int, float, bool, string,
także tablice i zbiory

zmienne decyzyjne – ustalane przez solver
(mogą mieć przypisaną wartość początkową),
dziedziny takie same, jak dla parametrów,
możliwość definiowania przedziałów

ograniczenia, jakie muszą spełniać zmienne
decyzyjne (rozwiązanie), dostępne typowe
operatory relacyjne oraz bardziej złożone
predykaty

wyświetlenie listy z wynikami (tylko jedno
polecenie w programie), możliwość prostego
formatowania (funkcje show, show_int oraz
show_float), sklejanie łańcuchów przez ++

Problem 1: kolorowanie mapy



```
Compiling aust.mzn
Running aust.mzn
wa=3 nt=2 sa=1 q=3 nsw=2 v=3 t=1
-----
wa=3 nt=2 sa=1 q=3 nsw=2 v=3 t=2
-----
wa=3 nt=2 sa=1 q=3 nsw=2 v=3 t=3
-----
wa=2 nt=3 sa=1 q=2 nsw=3 v=2 t=1
-----
wa=2 nt=3 sa=1 q=2 nsw=3 v=2 t=2
-----
wa=2 nt=3 sa=1 q=2 nsw=3 v=2 t=3
-----
wa=3 nt=1 sa=2 q=3 nsw=1 v=3 t=1
-----
wa=3 nt=1 sa=2 q=3 nsw=1 v=3 t=2
-----
wa=3 nt=1 sa=2 q=3 nsw=1 v=3 t=3
-----
wa=1 nt=3 sa=2 q=1 nsw=3 v=1 t=1
-----
wa=1 nt=3 sa=2 q=1 nsw=3 v=1 t=2
-----
wa=1 nt=3 sa=2 q=1 nsw=3 v=1 t=3
-----
wa=2 nt=1 sa=3 q=2 nsw=1 v=2 t=1
-----
wa=2 nt=1 sa=3 q=2 nsw=1 v=2 t=2
-----
wa=2 nt=1 sa=3 q=2 nsw=1 v=2 t=3
-----
wa=1 nt=2 sa=3 q=1 nsw=2 v=1 t=1
-----
wa=1 nt=2 sa=3 q=1 nsw=2 v=1 t=2
-----
wa=1 nt=2 sa=3 q=1 nsw=2 v=1 t=3
-----
=====
Finished in 155msec
```

wykorzystanie solvera *finite domain* do
znalezienia rozwiązań

Problem 2: planowanie wypieków



ograniczenia na zasoby:
mąka, banany, cukier, masło, kakao

```
1 var 0..100: b;  
2 var 0..100: c;  
3  
4 constraint 250*b + 200*c <= 4000;  
5 constraint 2*b <= 6;  
6 constraint 75*b + 150*c <= 2000;  
7 constraint 100*b + 150*c <= 500;  
8 constraint 75*c <= 500;  
9  
10 solve maximize 400*b + 450*c;  
11  
12 output ["# of banana cakes=\(b) # of chocolate cakes=\(c)"];
```

definicja funkcji celu: typowe operatory
arytmetyczne, div i mod, typowe funkcje
matematyczne

```
Compiling cakes.mzn  
Running cakes.mzn  
# of banana cakes=2 # of chocolate cakes=2  
-----  
=====  
Finished in 70msec
```

Zbiory danych i asercje

- Możliwość zewnętrznego ustawienia wartości parametrów oraz zmiennych decyzyjnych
- Zewnętrzny plik danych (.dzn) albo linia poleceń (opcja -D)
- Asercje służące do sprawdzania poprawności danych wejściowych (także do ustalania wartości)

Problem 2: planowanie wypieków (wydzielone dane)

```
1 flour = 4000;  
2 banana = 6;  
3 sugar = 2000;  
4 butter = 500;  
5 cocoa = 500;
```

polecenie output można opuścić → wyjście w formacie zgodnym z formatem pliku danych

```
1 flour = 8000;  
2 banana = 11;  
3 sugar = 3000;  
4 butter = 1500;  
5 cocoa = 800;
```

```
1 int: flour;  
2 int: banana;  
3 int: sugar;  
4 int: butter;  
5 int: cocoa;
```

```
7 constraint assert(flour >= 0, "Amount of flour should be non-negative");  
8 constraint assert(banana >= 0, "Amount of banana should be non-negative");  
9 constraint assert(sugar >= 0, "Amount of sugar should be non-negative");  
10 constraint assert(butter >= 0, "Amount of butter should be non-negative");  
11 constraint assert(cocoa >= 0, "Amount of cocoa should be non-negative");
```

```
13 var 0..100: b;  
14 var 0..100: c;  
16 constraint 250*b + 200*c <= flour;  
17 constraint 2*b <= banana;  
18 constraint 75*b + 150*c <= sugar;  
19 constraint 100*b + 150*c <= butter;  
20 constraint 75*c <= cocoa;
```

```
22 solve maximize 400*b + 450*c;  
23  
24 output ["# of banana cakes=\\(b)\\n# of chocolate cakes=\\(c)"];
```

Compiling cakes2.mzn, with additional data pantry1.dzn
Running cakes2.mzn
of banana cakes=2
of chocolate cakes=2

Finished in 104msec

Compiling cakes2.mzn, with additional data pantry2.dzn
Running cakes2.mzn
of banana cakes=3
of chocolate cakes=8

Finished in 103msec

Problem 3: spłata pożyczki



```
1 var float: R;  
2 var float: P;  
3 var 0.0..5.0: I;
```

```
4  
5 var float: B1;  
6 var float: B2;  
7 var float: B3;  
8 var float: B4;  
9  
10 constraint B1 = P*(1.0 + I/4) - R;  
11 constraint B2 = B1*(1.0 + I/4) - R;  
12 constraint B3 = B2*(1.0 + I/4) - R;  
13 constraint B4 = B3*(1.0 + I/4) - R;  
14  
15 solve satisfy;
```

```
16  
17 output [  
18   "Borrowing \(\theta, 2, P) at \((I*100.0)\% interest\n",  
19   "and repaying \(\theta, 2, R) per quarter leaves \(\theta, 2, B4) owing"  
20 ];
```

parametry i zmienne decyzyjne typu float –
możliwość ograniczenia wartości do przedziału

możliwość formatowania zmiennych decyzyjnych
osadzonych w łańcuchu znakowym (takie samo,
jak przy użyciu funkcji show_...)

Problem 3: spłata pożyczki

```
1 I = 0.16;  
2 P = 1000.0;  
3 R = 260.0;
```



```
15 solve satisfy;  
16  
17 output [  
18   "Borrowing \(\theta, 2, P) at \((I*100.0)\% interest\n",  
19   "and repaying \(\theta, 2, R) per quarter leaves \(\theta, 2, B4) owing"  
20 ];
```



```
Compiling loan.mzn with data loan1.dzn  
Running loan.mzn  
Borrowing 1000.00 at 16.0% interest  
and repaying 260.00 per quarter leaves 65.78 owing  
-----  
Finished in 168msec
```

```
1 I = 0.16;  
2 P = 1000.0;  
3 B4 = 0.0;
```



```
Compiling loan.mzn with data loan2.dzn  
Running loan.mzn  
Borrowing 1000.00 at 16.0% interest  
and repaying 275.49 per quarter leaves 0.00 owing  
-----  
Finished in 153msec
```

```
1 I = 0.16;  
2 R = 250.0;  
3 B4 = 0.0;
```



wykorzystanie solvera *mixed integer programming* do znalezienia rozwiązania

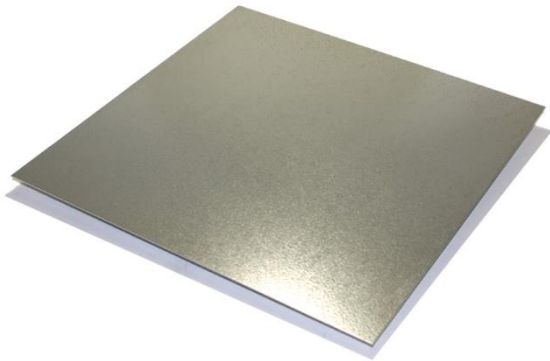


```
Compiling loan.mzn with data loan3.dzn  
Running loan.mzn  
Borrowing 907.47 at 16.0% interest  
and repaying 250.00 per quarter leaves 0.00 owing  
-----  
Finished in 153msec
```

Tablice i zbiory

- Tablice (wielowymiarowe) → złożone struktury z parametrami oraz zmiennymi decyzyjnymi
- Zbiory → definiowanie nowych typów danych, np. służących do indeksowania tablic
- Podstawowe operacje na zbiorach, np.: `in`, `subset`, `superset`, `union`, `intersect`, `diff`, `symdiff`, `card`

Problem 4: temperatura arkusza metalu



```
1 int: w = 4;  
2 int: h = 4;  
3  
4 set of int: HEIGHT = 0..h;  
5 set of int: IHEIGHT = 1..h-1;  
6 set of int: WIDTH = 0..w;  
7 set of int: IWIDTH = 1..w-1;  
8  
9 array[HEIGHT, WIDTH] of var float: t;
```

tablica zmiennych decyzyjnych

```
11 var float: left;  
12 var float: right;  
13 var float: top;  
14 var float: bottom;  
15  
16 constraint forall(i in IHEIGHT, j in IWIDTH)  
17   (4.0*t[i, j] = t[i-1, j] + t[i, j-1] + t[i+1, j] + t[i, j+1]);  
18  
19 constraint forall(i in IHEIGHT)(t[i, 0] = left);  
20 constraint forall(i in IHEIGHT)(t[i, w] = right);  
21 constraint forall(j in IWIDTH)(t[0, j] = top);  
22 constraint forall(j in IWIDTH)(t[h, j] = bottom);  
23  
24 constraint t[0,0] = 0.0;  
25 constraint t[0,w] = 0.0;  
26 constraint t[h,0] = 0.0;  
27 constraint t[h,w] = 0.0;  
28  
29 left = 0.0;  
30 right = 0.0;  
31 top = 100.0;  
32 bottom = 0.0;
```

metoda elementów skończonych,
równanie Laplace'a

Problem 4: temperatura arkusza metalu

```
34 solve satisfy;  
35  
36 output [show_float(6, 2, t[i, j]) ++  
37 if j == w then "\n" else " " endif |  
38 i in HEIGHT, j in WIDTH  
39 ];
```

wykorzystanie solvera *mixed integer programming* do znalezienia rozwiązania

Compiling laplace.mzn
Running laplace.mzn
-0.00 100.00 100.00 100.00 -0.00
-0.00 42.86 52.68 42.86 -0.00
-0.00 18.75 25.00 18.75 -0.00
-0.00 7.14 9.82 7.14 -0.00
-0.00 -0.00 -0.00 -0.00 -0.00

Finished in 419msec

Problem 5: planowanie produkcji

- Rozszerzenie problemu wypieku ciasta na dowolną listę produktów oraz zasobów (dane wejściowe)
- Rozdzielenie deklaracji i definicji typu wyliczeniowego
- Wykorzystanie typów wyliczeniowych do indeksowania tablic

Problem 5: planowanie produkcji

```
1 enum Products;
2 array[Products] of int: profit;
3
4 enum Resources;
5 array[Resources] of int: capacity;
6
7 array[Products, Resources] of int: consumption;
8 constraint assert(forall (p in Products, r in Resources)
9                 (consumption[p, r] >= 0), "Error: negative consumption");
10
11 int: mproducts = max (p in Products)
12                    (min (r in Resources where consumption[p, r] > 0)
13                      (capacity[r] div consumption[p, r]));
14
15 array[Products] of var 0..mproducts: produced;
16 array[Resources] of var 0..max(capacity): used;
17
18 constraint forall (r in Resources) (
19   used[r] = sum (p in Products)(consumption[p, r]*produced[p])
20   /\ used[r] <= capacity[r]
21 );
22
23 solve maximize sum (p in Products) (profit[p]*produced[p]);
24
25 output
26   ["\ (p) = \ (produced[p]);\n" | p in Products] ++
27   ["\ (r) = \ (used[r]);\n" | r in Resources];
```

tabele z parametrami, sprawdzenie poprawności danych wejściowych

oszacowanie maksymalnej liczby wyprodukowanych produktów

tabele ze zmiennymi decyzyjnymi i ograniczenia narzucone na zmienne

funkcja celu

Problem 5: planowanie produkcji

```
1 Products = {BananaCake, ChocolateCake};
2 profit = [400, 450];
3
4 Resources = {Flour, Banana, Sugar, Butter, Cocoa};
5 capacity = [4000, 6, 2000, 500, 500];
6
7 consumption = [ | 250, 2, 75, 100, 0,
8                 | 200, 0, 150, 150, 75 |];
```

definicje typów wyliczeniowych oraz
wartości parametrów

definicja tabeli dwuwymiarowej, także
funkcje `arraynd(indeksy, lista)`

```
23 solve maximize sum (p in Products) (profit[p]*produced[p]);
24
25 output
26 ["\ (p) = \ (produced[p]);\n" | p in Products] ++
27 ["\ (r) = \ (used[r]);\n" | r in Resources];
```

Sklejanie tabel jednowymiarowych (list) za
pomocą operatora ++

```
Compiling simple-prod-planning.mzn, with additional data simple-prod-planning-data.dzn
Running simple-prod-planning.mzn
BananaCake = 2;
ChocolateCake = 2;
Flour = 900;
Banana = 4;
Sugar = 450;
Butter = 500;
Cocoa = 150;
-----
=====
Finished in 311msec
```


Definiowanie list i zbiorów ($\rightarrow$ *List and set comprehension*)

[<expr> | <generator-expr>]



<generator>
<generator> where <bool-expr>



<identifier>, ..., <identifier> in <array-expr>

[i + j | i,j in 1..3 where i < j]



[1+2, 1+3, 2+3] $\rightarrow$ [3, 4, 5]

Wywoływanie generatorów (→ *Generator call expressions*)

```
<agg-func> (<generator-expr>) (<expr>)
```

```
<agg-func> ([<expr> | <generator-expr>])
```

Funkcje agregujące

- tablice liczbowe: `sum`, `product`, `min`, `max`
- tablice logiczne: `forall`, `exists`, `xorall`, `iffall`
(nieparzysta i parzysta liczba spełnionych warunków)

```
forall([a[i] != a[j] | i,j in 1..3 where i < j])
```

Problem 6: send more money...

Handwritten solution for the cryptarithm "SEND + MORE = MONEY":

$$\begin{array}{r} 9517 \\ + 1085 \\ \hline 10652 \end{array}$$

```
1 include "alldifferent.mzn";
2
3 var 1..9: S;
4 var 0..9: E;
5 var 0..9: N;
6 var 0..9: D;
7 var 1..9: M;
8 var 0..9: O;
9 var 0..9: R;
10 var 0..9: Y;
11
12 constraint 1000*S + 100*E + 10*N + D
13           + 1000*M + 100*O + 10*R + E
14           = 10000*M + 1000*O + 100*N + 10*E + Y;
15
16 constraint alldifferent([S,E,N,D,M,O,R,Y]);
17
18 solve satisfy;
19
20 output [ "  \S)\(E)\(N)\(D)\n",
21          "+  \M)\(O)\(R)\(E)\n",
22          "= \M)\(O)\(N)\(E)\(Y)\n"];
23
```

predefiniowany predykat sprawdzający
unikalność poszczególnych wartości

Compiling send-more-money.mzn
Running send-more-money.mzn
9567
+ 1085
= 10652

Finished in 350msec

Wyrażenia warunkowe

```
if <bool-expr> then <expr-1> else <expr-2> endif
```

```
int: r = if y != 0 then x div y else 0 endif;
```

```
constraint forall(i,j in PuzzleRange) (  
    if start[i,j] > 0  
    then puzzle[i,j] = start[i,j]  
    else true endif );
```

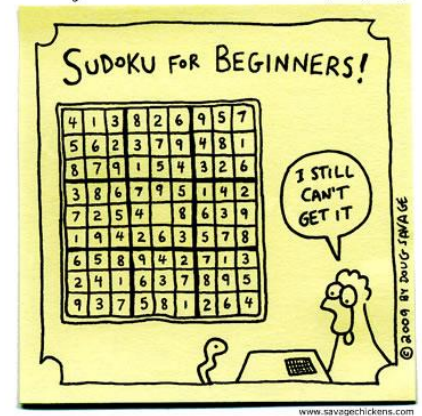
```
if j mod S == 0 then " " else "" endif
```

Problem 7: sudoku

```
1 include "alldifferent.mzn";
2
3 int: S;
4 int: N = S*S;
5 int: digs = ceil(log(10.0, int2float(N)));
6
7 set of int: PuzzleRange = 1..N;
8 set of int: SubSquareRange = 1..S;
9
10 array[PuzzleRange, PuzzleRange] of 0..N: start;
11 array[PuzzleRange, PuzzleRange] of var PuzzleRange: puzzle;
12
13 constraint forall(i, j in PuzzleRange) (
14   if start[i, j] > 0 then puzzle[i, j] = start[i, j] else true endif);
15
16 constraint forall (i in PuzzleRange)
17   (alldifferent([puzzle[i, j] | j in PuzzleRange]));
18 constraint forall (j in PuzzleRange)
19   (alldifferent([puzzle[i, j] | i in PuzzleRange]));
20
21 constraint forall (bi, bj in SubSquareRange)
22   (alldifferent([puzzle[(bi - 1)*S + i, (bj - 1)*S + j] | i, j in SubSquareRange]));
23
24 solve satisfy;
25
26 output [
27   show_int(digs, puzzle[i, j]) ++ " " ++
28   if j mod S == 0 then " " else "" endif ++
29   if j == N /\ i != N then
30     if i mod S == 0 then "\n\n" else "\n" endif
31   else "" endif
32   | i, j in PuzzleRange] ++ ["\n"];
```

Savage Chickens

by Doug Savage

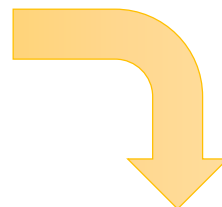


Problem 7: sudoku

```
1 S=3;
2 start=[|
3 0, 0, 0, 0, 0, 0, 0, 0, 0|
4 0, 6, 8, 4, 0, 1, 0, 7, 0|
5 0, 0, 0, 0, 8, 5, 0, 3, 0|
6 0, 2, 6, 8, 0, 9, 0, 4, 0|
7 0, 0, 7, 0, 0, 0, 9, 0, 0|
8 0, 5, 0, 1, 0, 6, 3, 2, 0|
9 0, 4, 0, 6, 1, 0, 0, 0, 0|
10 0, 3, 0, 2, 0, 7, 6, 9, 0|
11 0, 0, 0, 0, 0, 0, 0, 0, 0|];
```



```
23
24 solve satisfy;
25
```



```
Compiling sudoku.mzn, with additional data sudoku.dzn
Running sudoku.mzn
5 9 3 7 6 2 8 1 4
2 6 8 4 3 1 5 7 9
7 1 4 9 8 5 2 3 6

3 2 6 8 5 9 1 4 7
1 8 7 3 2 4 9 6 5
4 5 9 1 7 6 3 2 8

9 4 2 6 1 8 7 5 3
8 3 5 2 4 7 6 9 1
6 7 1 5 9 3 4 8 2
-----
=====
Finished in 261msec
```

Typy wyliczeniowe

```
enum <enum-name> = { ... } / anon_enum  
[var] <enum-name> : <var-name>  
[var] <l>..<u> : <var-name>
```

- Specjalizowane funkcje do obsługi enum
 - `enum_next(X, x)` – poprzednia wartość przed `x`
 - `enum_next(X, x)` – następna wartość po `x`
 - `to_enum(X, i)` – konwersja liczby `i` do wartości z `X`
- Standardowe funkcje stosowalne do enum
 - `card(X)` – liczba elementów w typie `X`
 - `min(X)` – pierwszy element w typie `X`
 - `max(X)` – ostatni element w typie `X`

Złożone ograniczenia

- Złożone wyrażenia logiczne w ograniczeniach
- Zastosowanie składni „matematycznej”
 - true, false
 - not
 - koniunkcja (/ \) i dysjunkcja (\ /)
 - *only-if* (< -), *implies* (- >), if-and-only-if (< - >)

Problem 8: warsztat (*jobshop*)



```
1 enum JOB;
2 enum TASK;
3 TASK: last = max(TASK);
4 array [JOB, TASK] of int: d;
5 int: total = sum(i in JOB, j in TASK)(d[i, j]);
6 int: digs = ceil(log(10.0, int2float(total)));
7 array [JOB, TASK] of var 0..total: s;
8 var 0..total: end;
9
10 constraint
11   forall(i in JOB) (
12     forall(j in TASK where j < last)
13       (s[i,j] + d[i,j] <= s[i, enum_next(TASK, j)])
14       /\ s[i,last] + d[i,last] <= end
15   );
16
17 constraint
18   forall(j in TASK) (
19     forall(i, k in JOB where i < k) (
20       s[i,j] + d[i,j] <= s[k,j] /\
21       s[k,j] + d[k,j] <= s[i,j]
22     )
23   );
24
25 % constraint end = 30;
26 solve minimize end;
27 % solve satisfy;
28
29 output ["end = \"(end)\\n\""] ++
30   [show_int(digs, s[i,j]) ++ " " ++
31     if j == last then "\\n" else "" endif | i in JOB, j in TASK];
32
```

zastosowanie typów enum jako indeksów
w tablicach (konwersja jednostronna)

Problem 8: warsztat (*jobshop*)

```
1 JOB = anon_enum(5);  
2 TASK = anon_enum(5);  
3 d = [ | 1, 4, 5, 3, 6  
4       | 3, 2, 7, 1, 2  
5       | 4, 4, 4, 4, 4  
6       | 1, 1, 1, 6, 8  
7       | 7, 3, 2, 2, 1 |];  
8
```



```
25 % constraint end = 30;  
26 solve minimize end;  
27 % solve satisfy;
```



```
-----  
end = 30  
6 9 13 18 21  
7 13 18 25 27  
1 5 9 13 17  
0 1 2 3 9  
12 22 25 27 29  
-----
```



```
25 constraint end = 30;  
26 % solve minimize end;  
27 solve satisfy;
```



```
-----  
end = 30  
6 9 13 18 21  
7 13 18 25 27  
1 5 9 13 17  
0 1 2 3 9  
13 22 25 27 29  
-----
```

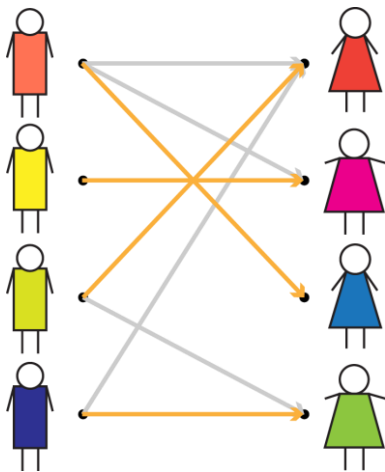


```
-----  
end = 30  
6 9 13 18 21  
7 13 18 25 27  
1 5 9 13 17  
0 1 2 3 9  
14 22 25 27 29  
-----
```

```
-----  
end = 30  
6 9 13 18 21  
7 13 18 25 27  
1 5 9 13 17  
0 1 2 3 9  
15 22 25 27 29  
-----
```

```
[ 100 more solutions ]  
[ 200 more solutions ]  
[ 400 more solutions ]  
[ 800 more solutions ]  
[ 1600 more solutions ]  
[ 3200 more solutions ]  
[ 6400 more solutions ]  
Stopped.  
[ 12800 more solutions ]  
Finished in 6s 717msec
```

Problem 9: planowanie stabilnych małżeństw



```
1 int: n;
2
3 enum Men = anon_enum(n);
4 enum Women = anon_enum(n);
5
6 array[Women, Men] of int: rankWomen;
7 array[Men, Women] of int: rankMen;
8
9 array[Men] of var Women: wife;
10 array[Women] of var Men: husband;
11
12 constraint forall (m in Men) (husband[wife[m]]=m);
13 constraint forall (w in Women) (wife[husband[w]]=w);
14
15 constraint forall (m in Men, o in Women) (
16   rankMen[m,o] < rankMen[m,wife[m]] -> rankWomen[o,husband[o]] < rankWomen[o,m]
17 );
18
19 constraint forall (w in Women, o in Men) (
20   rankWomen[w, o] < rankWomen[w, husband[w]] -> rankMen[o, wife[o]] < rankMen[o, w]
21 );
22
23 solve satisfy;
24
25 output ["wives=\\(wife)\\nhusbands=\\(husband)\\n"];
```

wymuszenie „stabilności” z punktu widzenia rankingu (preferencji)

Problem 9: planowanie stabilnych małżeństw

```
1 n = 5;
2 rankWomen =
3 [| 1, 2, 4, 3, 5,
4 | 3, 5, 1, 2, 4,
5 | 5, 4, 2, 1, 3,
6 | 1, 3, 5, 4, 2,
7 | 4, 2, 3, 5, 1 |];
8 rankMen =
9 [| 5, 1, 2, 4, 3,
10 | 4, 1, 3, 2, 5,
11 | 5, 3, 2, 4, 1,
12 | 1, 5, 4, 3, 2,
13 | 4, 3, 2, 1, 5 |];
```

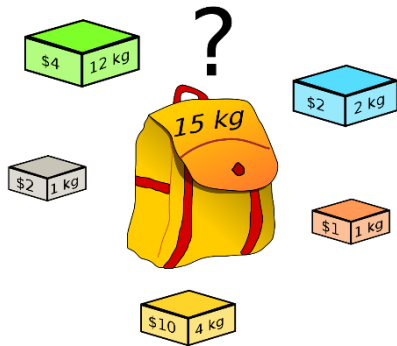


```
23 solve satisfy;
24
25 output ["wives=\(wife)\nhusbands=\(husband)\n"];
```



```
Compiling stable-marriage.mzn, with additional data stable-marriage.dzn
Running stable-marriage.mzn
wives=[Women_2, Women_1, Women_5, Women_3, Women_4]
husbands=[Men_2, Men_1, Men_4, Men_5, Men_3]
-----
wives=[Women_2, Women_3, Women_5, Women_1, Women_4]
husbands=[Men_4, Men_1, Men_2, Men_5, Men_3]
-----
wives=[Women_4, Women_1, Women_2, Women_3, Women_5]
husbands=[Men_2, Men_3, Men_4, Men_1, Men_5]
-----
=====
Finished in 212msec
```

Problem 10: problem plecakowy (*knapsack*)



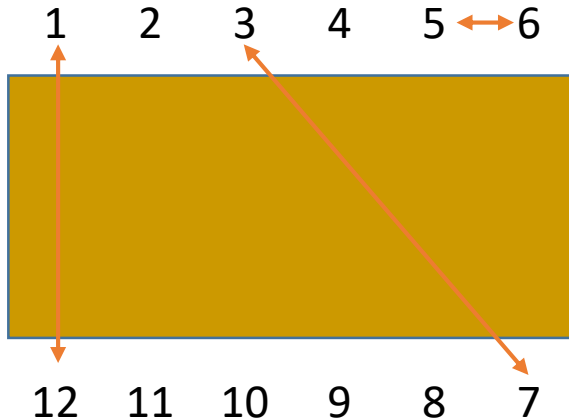
```
1 enum ITEM;  
2 int: capacity;  
3  
4 array[ITEM] of int: profits;  
5 array[ITEM] of int: weights;  
6  
7 var set of ITEM: knapsack;  
8  
9 constraint sum (i in knapsack) (weights[i]) <= capacity;  
10  
11 solve maximize sum (i in knapsack) (profits[i]);  
12  
13 output ["knapsack = \"(knapsack)\"\\n"];  
14
```

zbiór jako złożona zmienna decyzyjna

```
1 capacity = 16;  
2 ITEM = anon_enum(5);  
3 profits = [5, 15, 2, 6, 1];  
4 weights = [3, 10, 5, 3, 3];
```

```
> mzn-g12fd.bat knapsack.mzn knapsack.dzn  
knapsack = {ITEM_1, ITEM_2, ITEM_4}  
-----  
=====
```

Problem 11: planowanie miejsc przy stole weselnym



- Stół prostokątny, 12-miejscowy
- Wymagania ogólne (np. państwo młodzi razem)
- Antypatie i fobie gości
- Maksymalizacja odległości pomiędzy „antypatiami”

Problem 11: planowanie miejsc przy stole weselnym

```
1 enum Guests = { bride, groom, bestman, bridesmaid, bob, carol, ted, alice, ron, rona, ed, clara};
2 set of int: Seats = 1..12;
3 set of int: Hatreds = 1..5;
4 array[Hatreds] of Guests: h1 = [groom, carol, ed, bride, ted];
5 array[Hatreds] of Guests: h2 = [clara, bestman, ted, alice, ron];
6 set of Guests: Males = {groom, bestman, bob, ted, ron, ed};
7 set of Guests: Females = {bride, bridesmaid, carol, alice, rona, clara};
8
9 array[Guests] of var Seats: pos;
10 array[Hatreds] of var Seats: p1;
11 array[Hatreds] of var Seats: p2;
12 array[Hatreds] of var 0..1: sameside;
13 array[Hatreds] of var Seats: distance;
14
15 include "alldifferent.mzn";
16
17 constraint alldifferent(pos);
18 constraint forall (g in Males) (pos[g] mod 2 == 1);
19 constraint forall (g in Females) (pos[g] mod 2 == 0);
20
21 constraint not (pos[ed] in {1, 6, 7, 12});
22
23 constraint abs(pos[bride] - pos[groom]) <= 1
24 /\ (pos[bride] <= 6 <-> pos[groom] <= 6);
25
26 constraint forall(h in Hatreds) (
27     p1[h] = pos[h1[h]] /\
28     p2[h] = pos[h2[h]] /\
29     sameside[h] = bool2int(p1[h] <= 6 <-> p2[h] <= 6) /\
30     distance[h] = sameside[h] * abs(p1[h] - p2[h])
31     + (1 - sameside[h])*(abs(13 - p1[h] - p2[h]) + 1));
32
33 var int: total_distance;
34
35 constraint total_distance = sum(h in Hatreds)(distance[h]);
36
37 solve maximize total_distance;
38
39 output ["total distance = \"(total_distance)\\n\""] ++ [show(g) ++ " " | s in Seats, g in Guests where fix(pos[g]) == s] ++ ["\\n"];
```

antypatie wśród gości

panie siedzą na miejscach parzystych,
panowie na nieparzystych

Ed ma fobie – nie może siedzieć na miejscu
z brzegu stołu

państwo młodzi muszą siedzieć obok
siebie, po tej samej stronie stołu

maksymalizacja odległości pomiędzy
gośćmi, którzy za sobą nie przepadają

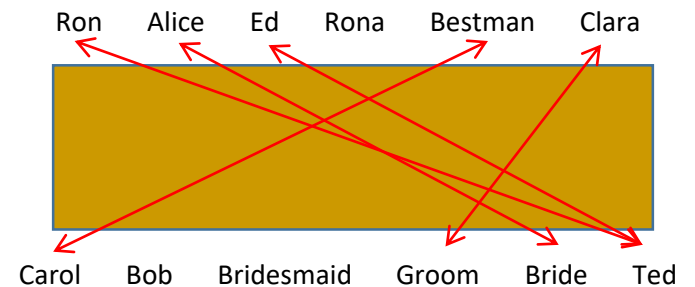
Problem 11: planowanie miejsc przy stole weselnym

```
solve maximize total_distance;

output ["total distance = \"(total_distance)\\n\""] ++
  [show(g) ++ " " | s in Seats, g in Guests where fix(pos[g]) == s] ++ ["\\n"];
```



```
Compiling wedding.mzn
Running wedding.mzn
total distance = 18
ron clara ed alice bestman rona ted bride groom carol bob bridesmaid
-----
total distance = 20
ron alice ed rona bestman clara ted bride groom carol bob bridesmaid
-----
total distance = 22
ron alice ed rona bestman clara ted bride groom bridesmaid bob carol
-----
=====
Finished in 834msec
```



Definiowanie własnych predykatów

```
predicate <pred-name>(<arg-def>,...,<arg-def>) = <bool-expr>
```

```
predicate lookup(array[int] of var int:x, int: i, var int: y) =  
    assert(i in index_set(x),  
        "index out of range in lookup"  
        y = x[i]);
```

```
predicate no_overlap(var int:s1, int:d1, var int:s2, int:d2) =  
    s1 + d1 <= s2 /\ s2 + d2 <= s1;
```

Definiowanie własnych predykatów

jobshop

```
16
17 constraint
18   forall(j in TASK) (
19     forall(i, k in JOB where i < k) (
20       s[i,j] + d[i,j] <= s[k,j] \/
21       s[k,j] + d[k,j] <= s[i,j]
22     )
23   );
24
```



```
17 predicate no_overlap(var int:s1, int:d1, var int:s2, int:d2) =
18   s1 + d1 <= s2 \/ s2 + d2 <= s1;
19
20 constraint
21   forall(j in TASK) (
22     forall(i, k in JOB where i < k) (
23       no_overlap(s[i,j], d[i,j], s[k, j], d[k,j])
24     )
25   );

```

Definiowanie własnych funkcji

```
function <ret-type> <func-name>(<arg-def>, ..., <arg-def>) = <expr>
```

```
function var int: manhattan(  
    var int: x1, var int: y1, var int: x2, var int: y2) =  
    abs(x1 - x2) + abs(y1 - y2);
```

```
function int: posn(int: a, int: a1) = (a-1)*S + a1;
```

Definiowanie własnych funkcji

sudoku

```
20  
21 constraint forall (bi, bj in SubSquareRange)  
22   (alldifferent([puzzle[(bi - 1)*S + i, (bj - 1)*S + j] | i, j in SubSquareRange]));  
23
```



```
21 function int: posn(int: a, int: a1) = (a-1)*S + a1;  
22  
23 constraint forall (bi, bj in SubSquareRange)  
24   (alldifferent([puzzle[posn(bi, i), posn(bj, j)] | i, j in SubSquareRange]));
```

Zmienne lokalne

```
let {<var-def>, ..., <var-def>} in <expr>
```

```
constraint let {var int: s = x1 + x2 + x3 + x4}  
            in l <= s /\ s <= u;
```

Zmienne lokalne

wedding

```
9 array[Guests] of var Seats: pos;  
10 array[Hatreds] of var Seats: p1;  
11 array[Hatreds] of var Seats: p2;  
12 array[Hatreds] of var 0..1: sameside;  
13 array[Hatreds] of var Seats: distance;
```

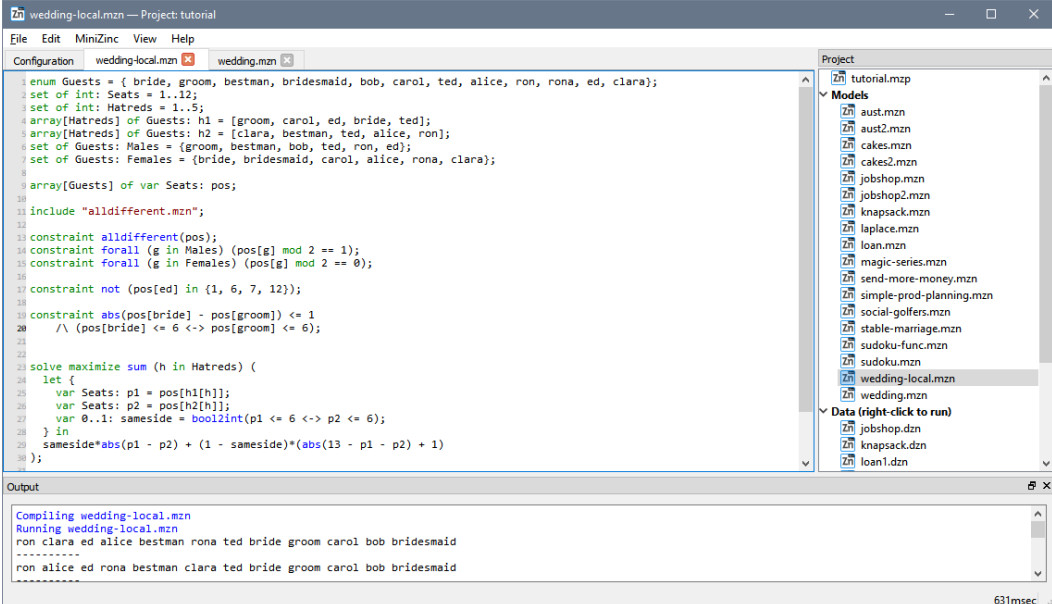
```
26 constraint forall(h in Hatreds) (  
27   p1[h] = pos[h1[h]] /\  
28   p2[h] = pos[h2[h]] /\  
29   sameside[h] = bool2int(p1[h] <= 6 <-> p2[h] <= 6) /\  
30   distance[h] = sameside[h] * abs(p1[h] - p2[h])  
31   + (1 - sameside[h])*(abs(13 - p1[h] - p2[h]) + 1));
```



```
9 array[Guests] of var Seats: pos;  
  
23 solve maximize sum (h in Hatreds) (  
24   let {  
25     var Seats: p1 = pos[h1[h]];  
26     var Seats: p2 = pos[h2[h]];  
27     var 0..1: sameside = bool2int(p1 <= 6 <-> p2 <= 6);  
28   } in  
29   sameside*abs(p1 - p2) + (1 - sameside)*(abs(13 - p1 - p2) + 1)  
30 );
```

Wykorzystanie „samodzielne”

- Proste środowisko do edycji plików (.mzn, .dzn) i uruchamiania obliczeń
- Zestaw narzędzi do uruchamiania z linii poleceń



```
enum Guests = { bride, groom, bestman, bridesmaid, bob, carol, ted, alice, ron, rona, ed, clara};
set of int: Seats = 1..12;
set of int: Hatsreds = 1..5;
array[Hatsreds] of Guests: h1 = [groom, carol, ed, bride, ted];
array[Hatsreds] of Guests: h2 = [clara, bestman, ted, alice, ron];
set of Guests: Males = {groom, bestman, bob, ted, ron, ed};
set of Guests: Females = {bride, bridesmaid, carol, alice, rona, clara};
array[Guests] of var Seats: pos;
include "alldifferent.mzn";
constraint alldifferent(pos);
constraint forall (g in Males) (pos[g] mod 2 == 1);
constraint forall (g in Females) (pos[g] mod 2 == 0);
constraint not (pos[ed] in {1, 6, 7, 12});
constraint abs(pos[bride] - pos[groom]) <= 1
/\ (pos[bride] <= 6 <-> pos[groom] <= 6);
solve maximize sum (h in Hatsreds) (
let {
var Seats: p1 = pos[h1[h]];
var Seats: p2 = pos[h2[h]];
var 0..1: sameside = bool2int(p1 <= 6 <-> p2 <= 6);
} in
sameside*abs(p1 - p2) + (1 - sameside)*(abs(13 - p1 - p2) + 1)
);
```

```
Compiling wedding-local.mzn
Running wedding-local.mzn
ron clara ed alice bestman rona ted bride groom carol bob bridesmaid
-----
ron alice ed rona bestman clara ted bride groom carol bob bridesmaid
```

631msec

Integracja z własnym oprogramowaniem

- Uruchomienie referencyjnych narzędzi z poziomu Python-a
- Możliwość przekazania danych z poziomu aplikacji oraz przejęcia wyników („surowe” i sformatowane)

Integracja z własnym oprogramowaniem

```
In [1]: import pymzn
```

```
In [2]: s = pymzn.minizinc('loan.mzn', 'loan1.dzn', solver=pymzn.g12mip)
print(s)
```

```
[{'B1': 780.0, 'B2': 551.2, 'B3': 313.248, 'B4': 65.77792000000005}]
```

```
In [3]: for sol in s:
        print("B1 = {:.2f}, B2 = {:.2f}, B3 = {:.2f}, B4 = {:.2f}".format(sol['B1'], sol['B2'], sol['B3'], sol['B4']))
```

```
B1 = 780.00, B2 = 551.20, B3 = 313.25, B4 = 65.78
```

```
In [4]: s = pymzn.minizinc('loan.mzn', 'loan1.dzn', solver=pymzn.g12mip, output_mode='item')
print(s)
```

```
['Borrowing 1000.00 at 16.0% interest\nand repaying 260.00 per quarter leaves 65.78 owing']
```

```
In [5]: s = pymzn.minizinc('loan.mzn', data={'I':0.16, 'P':1000.0, 'B4': 0.0}, solver=pymzn.g12mip, output_mode='dict')
print(s)
```

```
[{'R': 275.4900453648024, 'B1': 764.5099546351977, 'B2': 519.6003074558032, 'B3': 264.894274389233}]
```

```
In [6]: s = pymzn.minizinc('loan.mzn', data={'I':0.16, 'P':1000.0, 'B4': 0.0}, solver=pymzn.g12mip, output_mode='item')
print(s)
```

```
['Borrowing 1000.00 at 16.0% interest\nand repaying 275.49 per quarter leaves 0.00 owing']
```

Podsumowanie

- MiniZinc – język do (czytelnego) modelowania problemów decyzyjnych i optymalizacyjnych
- Dostępne solvery referencyjne, możliwa integracja z wieloma istniejącymi solverami
- Możliwość integracji z własnymi aplikacjami (Python, Java, ...)