

Towards Schema-independent Querying on Document Data Stores

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Documement-oriented Database

- **Data format:** Semi-structured documents, JSON, BSON ...
- **Data model:** Schema-less
- **Advantage:** Big data support, Scalability, Availability
- **Example:** MongoDB, CouchDB
- **Applications:** Web, IoT, social media ...
- **Interrogation:** JDBC, Drivers, API, Command line ...

```
{
  "firstName": "John",
  "lastName": "Smith",
  "age": 25,
  "address": {
    "streetAddress": "21 2nd Street",
    "city": "New York",
    "state": "NY",
    "postalCode": "10021"
  },
  "phoneNumber": [
    { "type": "home", "number": "212 555-1234" },
    { "type": "fax", "number": "646 555-4567" }
  ]
}
```



Modeling Multi-structured Data

Collection

$$C = \{d_1, \dots, d_c\}$$

Document

$$d_i = (k_i, v_i)$$

- k_i is the document' identifier.
- $v_i = \{a_{i,1} : v_{i,1}, \dots, a_{i,n} : v_{i,n_i}\}$ is the document' value.

Document Schema

$$s_i = \{p_1, \dots, p_m\}$$

where p_i is a path leading to leaf node in document d_i .

Collection Schema

$$S = \bigcup_{i=1}^{|C|} s_i$$

Structural Heterogeneity

Document 1

```
{
  "_id": 1,
  "title": "Fast and furious",
  "year": 2017,
  "language": "English"
}
```

Document 2

```
{
  "_id": 2,
  "title": "Titanic",
  "details": {
    "year": 1997,
    "language": "English"
  }
}
```

Document 3

```
{
  "_id": 3,
  "title": "Despicable Me 3",
  "year": 2017
}
```

Document 4

```
{
  "_id": 4,
  "title": "The Hobbit",
  "versions": [
    {
      "year": 2012,
      "language": "English"
    },
    {
      "year": 2013,
      "language": "French"
    }
  ]
}
```

Query Operators

Kernel of Unary Operators

$$k = \{\pi, \sigma\}$$

Projection Operator

$$\pi_{(A)}(C_{in}) = C_{out}$$

The project operator reduces the initial schemas of documents to a finite subset of attributes A .

Selection Operator

$$\sigma_{(P)}(C_{in}) = C_{out}$$

The select operator retrieves only documents that match the selection condition P expressed in normal form ($Norm_p$).

Querying Multi-structured Data Problem

$$\pi(\text{"title", "year"})(C)$$

Document 3

```
{
  "_id": 3,
  "title": "Despicable Me 3",
  "year": 2017
}
```

Document 4

```
{
  "_id": 4,
  "title": "The Hobbit",
  "versions": [
    {
      "year": 2012,
      "language": "English"
    },
    {
      "year": 2013,
      "language": "French"
    }
  ]
}
```

Document 1

```
{
  "_id": 1,
  "title": "Fast and furious",
  "year": 2017,
  "language": "English"
}
```

Document 2

```
"_id": 2,
"title": "Titanic",
"details": {
  "year": 1997,
  "language": "English"
}
```

Querying Multi-structured Data Problem

$$\pi(\text{"title" , "year"})(C)$$

Document 3

Document 1

```
{
  "_id": 1,
  "title": "Fast and furious",
  "year":2017
  "language": "English"
}
```

```
{
  "_id": 3,
  "title": "Despicable Me 3",
  "year":2017
}
```

Document 2

```
{
  "_id": 2,
  "title": "Titanic",
  "details":
    {
      "year":1997
      "language": "English"
    }
}
```

Document 4

```
{
  "_id": 4,
  "title": "The Hobbit",
  "versions":
    [
      {
        "year":2012
        "language": "English"
      },
      {
        "year":2013
        "language": "French"
      }
    ]
}
```

Querying Multi-structured Data Problem

$$\pi(\text{"title" , "year" , "details.year" , "versions.1.year" , "versions.2.year"})(C)$$

Querying Multi-structured Data Problem

$$\pi(\text{"title", "year", "details.year", "versions.1.year", "versions.2.year"})(C)$$

Document 1

```
{
  "_id": 1,
  "title": "Fast and furious",
  "year": 2017
  "language": "English"
}
```

Document 3

```
{
  "_id": 3,
  "title": "Despicable Me 3",
  "year": 2017
}
```

Document 2

```
{
  "_id": 2,
  "title": "Titanic",
  "details": {
    "year": 1997
    "language": "English"
  }
}
```

Document 4

```
{
  "_id": 4,
  "title": "The Hobbit",
  "versions": [
    {
      "year": 2012
      "language": "English"
    },
    {
      "year": 2013
      "language": "French"
    }
  ]
}
```



Plan

- 1 Introduction
- 2 Querying Heterogeneous Documents
- 3 Experiments
- 4 Conclusion & perspectives



Physical data transformation

- Flattening data.
- Using additional databases.
- Introducing new structures.

[(Chasseuretal., 2013), (Taharaetal., 2014)(Taharaetal., 2014)]

⇒ *Need to learn new schema.*

⇒ *Loss of initial document schemas/structures.*

⇒ *Need to re – build new schemas when structres are changed.*

Virtual data transformation

- Inferring existing schemas.
- Building an unified schema.
- Tracking different schemas versions.

[(Baazizi et al., 2017),(Ruiz et al., 2015),(Wang et al., 2015)]

⇒ *Need to learn new structures.*

⇒ *Querying is only limited to structural level.*

⇒ *Heterogeneity is manually managed to formulate application queries.*

EasyQ

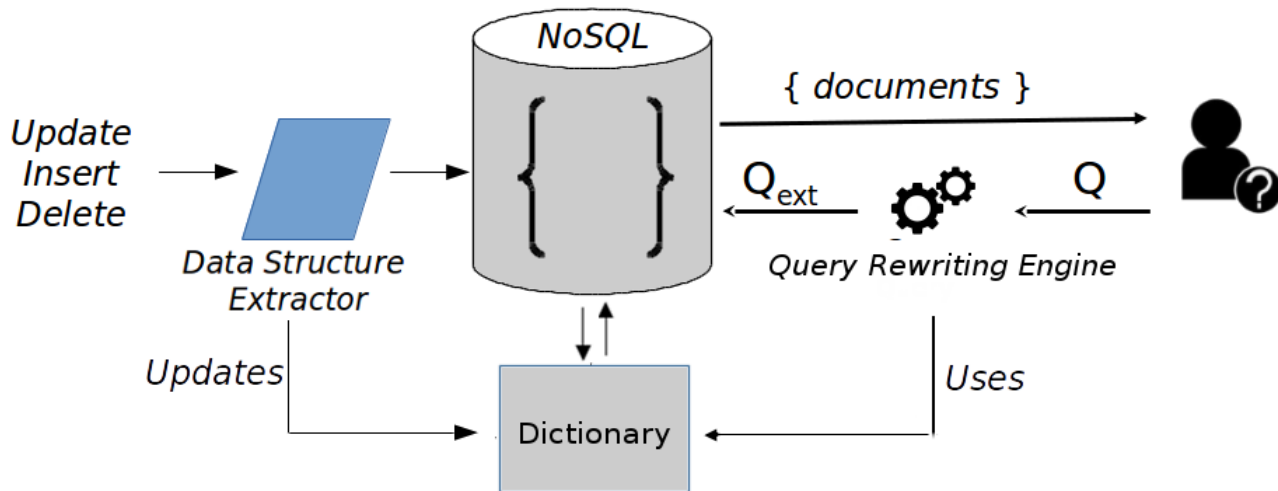


Figure: EasyQ Architecture

Dictionary

The dictionary $dict_C$ constructed from a collection C is defined by

$$dict_C = \{(p_k, \Delta_k)\} \forall p_k \in S_C$$

- $p_k \in S_C$ is a path leading to a leaf node which is present in at least one document;
- $\Delta_k = \{p_{p_k,1}, \dots, p_{p_k,q}\} \subseteq S_C$, is a set of navigational paths leading to p_k ;

Dictionary Construction Process

“year”

Document 1

```
{
  "_id": 1,
  "title": "Fast and furious",
  "year":2017,
  "language":"English"
}
```

Document 3

```
{
  "_id": 3,
  "title": "Despicable Me 3",
  "year":2017
}
```

Document 2

```
"_id": 2,
"title": "Titanic",
"details":
{
  "year":1997,
  "language":"English"
}
```

Document 4

```
{
  "_id": 4,
  "title": "The Hobbit",
  "versions":
  [{
    "year":2012,
    "language":"English"
  },
  {
    "year":2013
    "language":"French"
  }]
}
```



Dictionary Construction Process

$$dict = \{ (\text{“year”}, \{ \text{“year”}, \text{“details.year”}, \text{“versions.1.year”}, \text{“versions.2.year”} \}) \}$$

Dictionary

```
dict = {
  ("title", {"title"} ),
  ( "year", {"year", "details.year", "versions.1.year", "versions.2.year"}),
  ( "language", {"language", "details.language", "versions.1.language", "versions.2.language" } ),
  ( "details", {"details"} ),
  ( "details.year", {"details.year"} ),
  ( "details.language", {"details.language"}),
  ( "versions", {"versions"} ),
  ( "versions.1", {"version.1"} ) ,
  ( "versions.1.year", {"versions.1.year"} ),
  ( "versions.1.language", {"versions.1.language"} ),
  ( "versions.2", {"versions.2"} ),
  ( "versions.2.year", {"versions.2.year"} ),
  ( "versions.2.language", {"versions.2.language"} )
}
```

Algorithm for Automatic Query Extension

Algorithm 1: Automatic extension for the original user's query

Input: Q

Output: Q_{ext}

```

 $Q_{ext} \leftarrow id$  // identity
foreach  $q_i \in Q$  do
  switch  $q_i$  do
    case  $\pi_{A_i}$  // Projection operation
    do
       $A_{ext} \leftarrow \bigcup_{a_k \in A_i} \Delta_k$ ;  $Q_{ext} \leftarrow Q_{ext} \circ \pi_{A_{ext}}$ 
    end
    case  $\sigma_{Norm_p}$  //  $Norm_p = \bigwedge_i (\bigvee_j a_{i,j} \varpi_{i,j} v_{i,j})$  selection operation
    do
       $P_{ext} \leftarrow \bigwedge_k \left( \bigvee_l \bigvee_{a_j \in \Delta_{k,l}} a_j \varpi_{k,l} v_{k,l} \right)$ ;  $Q_{ext} \leftarrow Q_{ext} \circ \sigma_{P_{ext}}$ 
    end
  end
end
end
```

Extending project operator

Algorithm 1: Automatic extension for the original user's query

```

Input:  $Q$ 
Output:  $Q_{ext}$ 
 $Q_{ext} \leftarrow id$  // identity
foreach  $q_i \in Q$  do
  switch  $q_i$  do
    case  $\pi_{A_i}$  // Projection operation
    do
       $A_{ext} \leftarrow \bigcup_{a_k \in A_i} \Delta_k$ ;  $Q_{ext} \leftarrow Q_{ext} \circ \pi_{A_{ext}}$ 
    end
    case  $\sigma_{Norm_p}$  //  $Norm_p = \wedge_i (\vee_j a_{i,j} \varpi_{i,j} v_{i,j})$  selection operation
    do
       $P_{ext} \leftarrow \wedge_k \left( \vee_l \vee_{a_j \in \Delta_{k,l}} a_j \varpi_{k,l} v_{k,l} \right)$ ;  $Q_{ext} \leftarrow Q_{ext} \circ \sigma_{P_{ext}}$ 
    end
  end
end
end

```

Extending project operator

Attributes extensions

$$A_{ext} \leftarrow \bigcup_{a_k \in A_i} \Delta_k$$

Example

$$\pi(\text{"title"}, \text{"year"})(C)$$

Extending project operator

$$\pi(\text{"title"}, \text{"year"})(C)$$

```
dict = {
  ("title", {"title"} ),
  ("year", {"year", "details.year", "versions.1.year", "versions.2.year"})
  ("language", {"language", "details.language", "versions.1.language", "versions.2.language"}),
  ("details", {"details"} ),
  ("details.year", {"details.year"}),
  ("details.language", {"details.language"}),
  ("versions", {"versions"}),
  ("versions.1", {"version.1"}),
  ("versions.1.year", {"versions.1.year"} ),
  ("versions.1.language", {"versions.1.language"} ),
  ("versions.2", {"versions.2"} ),
  ("versions.2.year", {"versions.2.year"}),
  ("versions.2.language", {"versions.2.language"} )
}
```

Extending project operator

Attributes extensions

$$A_{ext} \leftarrow \bigcup_{a_k \in A_i} \Delta_k$$

Example

$$\pi(\text{"title"}, \text{"year"})(C)$$

- $A_{ext} \leftarrow \{\text{"title"}\} \cup \{\text{"year"}, \text{"details.year"}, \text{"versions.1.year"}, \text{"versions.2.year"}\}$

Projection query extended

$$\Rightarrow \pi(\text{"title"}, \text{"year"}, \text{"details.year"}, \text{"versions.1.year"}, \text{"versions.2.year"})(C)$$

Extending Select Operator

Algorithm 1: Automatic extension for the original user's query

```

Input:  $Q$ 
Output:  $Q_{ext}$ 
 $Q_{ext} \leftarrow id$  // identity
foreach  $q_i \in Q$  do
  switch  $q_i$  do
    case  $\pi_{A_i}$  // Projection operation
    do
       $A_{ext} \leftarrow \bigcup_{a_k \in A_i} \Delta_k$ ;  $Q_{ext} \leftarrow Q_{ext} \circ \pi_{A_{ext}}$ 
    end
    case  $\sigma_{Norm_p}$  //  $Norm_p = \bigwedge_i (\bigvee_j a_{i,j} \varpi_{i,j} v_{i,j})$  selection operation
    do
       $P_{ext} \leftarrow \bigwedge_k \left( \bigvee_l \bigvee_{a_j \in \Delta_{k,l}} a_j \varpi_{k,l} v_{k,l} \right)$ ;  $Q_{ext} \leftarrow Q_{ext} \circ \sigma_{P_{ext}}$ 
    end
  end
end

```

Extending Select Operator

Attributes extensions

$$P_{ext} \leftarrow \bigwedge_k \left(\bigvee_l \bigvee_{a_j \in \Delta_{k,l}} a_j \varpi_{k,l} v_{k,l} \right)$$

Example

$$\sigma("title" \neq \text{Null} \wedge "language" = \text{"English"})(C)$$

Extending Select Operator

Extending Selection's Predicates

$$P_{ext} \leftarrow \bigwedge_k \left(\bigvee_l \bigvee_{a_j \in \Delta_{k,l}} a_j \varpi_{k,l} v_{k,l} \right)$$

Example

$$\sigma(\text{"title"} \neq \text{Null} \wedge \text{"language"} = \text{"English"})(C)$$

Selection query extended

$$P_{ext} \leftarrow \left(\bigvee_{a_j \in \Delta(\text{"title"})} a_j \neq \text{Null} \right) \wedge \left(\bigvee_{a_j \in \Delta(\text{"language"})} a_j = \text{"English"} \right) \\ \Rightarrow \sigma(P_{ext})(C)$$

Extending Select Operator

$$\sigma\left(\bigvee_{a_j \in \Delta(\text{"title"})} a_j \neq \text{Null}\right) \wedge \left(\bigvee_{a_j \in \Delta(\text{"language"})} a_j = \text{"English"}\right)(C)$$

```
dict = {
  ("title", {"title"} ),
  ("year", {"year", "details.year", "versions.1.year", "versions.2.year"}),
  ("language", {"language", "details.language", "versions.1.language", "versions.2.language"}),
  ("details", {"details"} ),
  ("details.year", {"details.year"}),
  ("details.language", {"details.language"}),
  ("versions", {"versions"}),
  ("versions.1", {"version.1"} ),
  ("versions.1.year", {"versions.1.year"} ),
  ("versions.1.language", {"versions.1.language"} ),
  ("versions.2", {"versions.2"} ),
  ("versions.2.year", {"versions.2.year"}),
  ("versions.2.language", {"versions.2.language"} )
}
```

Extending Select Operator

Selection query extended

$$\sigma(\bigvee_{a_j \in \Delta(\text{"title"})} a_j \neq \text{Null}) \wedge (\bigvee_{a_j \in \Delta(\text{"language"})} a_j = \text{"English"}) (C)$$

Rewritten Query

$$\sigma \left(\left(\text{"title"} \neq \text{Null} \right) \wedge \left(\text{"language"} = \text{"English"} \vee \text{"details.language"} = \text{"English"} \vee \text{"versions.1.language"} = \text{"English"} \vee \text{"versions.2.language"} = \text{"English"} \right) (C)$$

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Synthetic dataset

```
{
  "_id" : "012cde54e24fa175a1ce7965d",
  "director_name" : "Bob Odenkirk",
  "gross" : 1954202,
  "movie_title" : "A Nightmare on Elm Street 4: The Dream Master ",
  "facenumber_in_poster" : 43,
  "plot_keywords" : "apartment|oven|stove|thanksgiving|thanksgiving dinner",
  "language" : "Mandarin",
  "title_year" : 1982,
  "aspect_ratio" : 4,
  "director_facebook_likes" : 177,
  "cast_total_facebook_likes" : 13716,
  "num_critic_for_reviews" : 478,
  "actor_2_name" : "George Sanders",
  "actor_1_facebook_likes" : 44000,
  "num_user_for_reviews" : 1416,
  "country" : "Philippines",
  "content_rating" : "TV-G",
  "actor_2_facebook_likes" : 162,
  "duration" : 201,
  "actor_3_name" : "Jason Bateman",
  "budget" : 26000000,
  "imdb_score" : 3.6,
  "actor_3_facebook_likes" : 794,
  "actor_1_name" : "Numan Acar",
  "movie_imdb_link" : "http://www.imdb.com/title/tt2622294/?ref_=fn_tt_tt_1"
}
```

Figure: Flat Document d1 Describing Movies from IMDB

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Synthetic dataset

```
{ "_id" : "59cdelle25fa005a1ce7966d",
  "group_18" : {
    "level0" : {
      "level1" : { "director_name" : "Bob Odenkirk",
                  "gross" : 1954202 }
    }
  },
  "group_28" : {
    "level0" : {
      "level1" : { "director_facebook_likes" : 177,
                  "cast_total_facebook_likes" : 13716 }
    }
  },
  "group_38" : {
    "level0" : {
      "level1" : { "movie_title" : "A Nightmare on Elm Street 4: The Dream Master ",
                  "facenumber_in_poster" : 43,
                  "plot_keywords" : "apartment|oven|stove|thanksgiving|thanksgiving dinner",
                  "language" : "Mandarin",
                  "title_year" : 1982,
                  "aspect_ratio" : 4 }
    }
  },
  "group_48" : {
    "level0" : {
      "level1" : { "num_critic_for_reviews" : 478,
                  "actor_2_name" : "George Sanders",
                  "actor_1_facebook_likes" : 44000,
                  "num_user_for_reviews" : 1416,
                  "country" : "Philippines",
                  "content_rating" : "TV-G",
                  "actor_2_facebook_likes" : 162 }
    }
  },
  "group_58" : {
    "level0" : {
      "level1" : { "duration" : 201,
                  "actor_3_name" : "Jason Bateman",
                  "budget" : 26000000,
                  "imdb_score" : 3.6 }
    }
  },
  "group_68" : {
    "level0" : {
      "level1" : { "actor_3_facebook_likes" : 794,
                  "actor_1_name" : "Numan Acar",
                  "movie_imdb_link" : "http://www.imdb.com/title/tt2622294/?ref_=fn_tt_tt_1" }
    }
  }
}
```

Figure: Document D1 after structural heterogeneity injection

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Settings of the generated dataset

Setting	Value
# of schema	10
# of grouping objects per schema	{5,6,1,3,4,2,7,2,1,3}
Nesting levels per schema	{4,2,6,1,5,7,2,8,3,4}
Percentage of schema presence	10%
# of attributes per schema	Random
# of attributes per grouping objects	Random
Collection size	10 GB, 25 GB, 50 GB, 100 GB
Number of documents per collection	12 M, 30 M , 60 M, 120 M

Table: Settings of the generated dataset

Queries predicates

Predicate	Attribute	Type	Operator	Paths	Depths	selectivity
p1	DirectorName	String	Regex{^A}	8	{8,2,3,9,6,5,4,7}	0,06 %
p2	Gross	Int	> 100 k	7	{7,8,2,3,9,6,4}	66 %
p3	Language	String	= "English"	7	{7,8,3,9,6,5,4}	0,018%
p4	Imdb_score	Float	<4,7	8	{8,7,2,3,4,5,6,9}	29 %
p5	Duration	Int	≤ 200	7	{7,8,2,3,6,5,4}	77%
p6	Country	String	≠ Null	6	{7,2,3,9,5,4}	100 %
p7	year	Int	< 1950	7	{7,8,2,3,6,5,4}	23 %
p8	FB_likes	Int	≥ 500	7	{6,2,3,8,5,4,3}	83 %

Table: Query predicates

Queries

Q1/Q2

$$\pi_{(*)}(\sigma(\text{director_name}="A\%" \ (\wedge/\vee)\text{gross}>100000)(C))$$

Q3/Q4

$$\pi_{(*)}(\sigma(\text{director_name}="A\%" \ (\wedge/\vee)\text{gross}>100000(\wedge/\vee)\text{duration}<200(\wedge/\vee)\text{title_year}<1950)(C))$$

Q5/Q6

$$\pi_{(*)}(\sigma(\text{director_name}="A\%" \ (\wedge/\vee)\text{gross}>100000(\wedge/\vee)\text{duration}<200(\wedge/\vee)\text{title_year}<1950$$

$$(\wedge/\vee)\text{pays}\neq\text{Null}(\wedge/\vee)\text{language}=\text{English}(\wedge/\vee)\text{imdb_score}<4(\wedge/\vee)$$

$$\text{cast_total_facebook_likes}>500)(C))$$


Queries

```
db.C.find(
  {'$and': [{'director_name': {'$regex': '^A'}}, {'gross': {'$gt': 100000}}]})
```



Queries

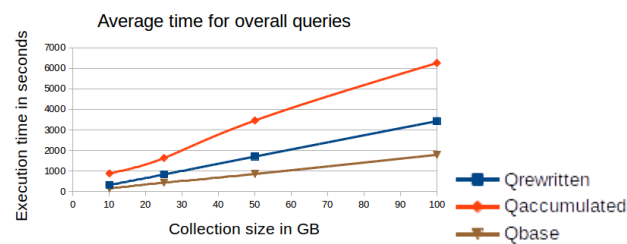
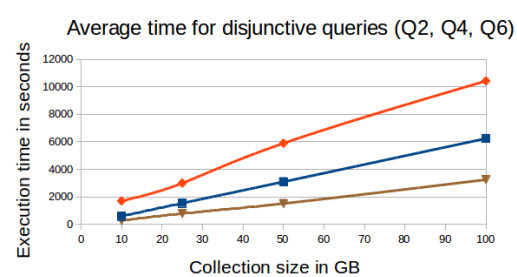
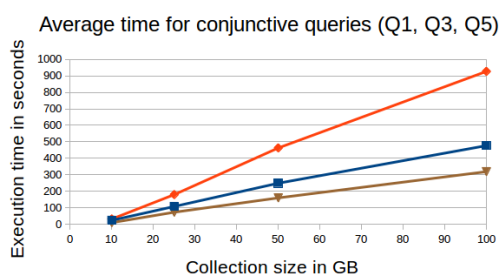
```

db.C.find(
  {'$and': [
    {'$or': [
      {u'group_1C.level0.level1.level2.level3.level4.level5.director_name': {'$regex': '^A'}},
      {u'group_1F.level0.level1.level2.level3.level4.level5.level6.director_name': {'$regex': '^A'}},
      {u'group_1D.level0.director_name': {'$regex': '^A'}}, {u'group_1B.level0.level1.director_name': {'$regex': '^A'}},
      {u'group_1H.level0.level1.level2.level3.level4.level5.level6.level7.director_name': {'$regex': '^A'}},
      {u'group_1E.level0.level1.level2.level3.level4.director_name': {'$regex': '^A'}},
      {u'group_1J.level0.level1.level2.level3.director_name': {'$regex': '^A'}},
      {u'group_1I.level0.level1.level2.director_name': {'$regex': '^A'}}]},
    [{u'group_1C.level0.level1.level2.level3.level4.level5.gross': {'$gt': 100000}},
     {u'group_1F.level0.level1.level2.level3.level4.level5.level6.gross': {'$gt': 100000}},
     {u'group_2D.level0.gross': {'$gt': 100000}},
     {u'group_1B.level0.level1.gross': {'$gt': 100000}},
     {u'group_2H.level0.level1.level2.level3.level4.level5.level6.level7.gross': {'$gt': 100000}},
     {u'group_3E.level0.level1.level2.level3.level4.gross': {'$gt': 100000}},
     {u'group_1I.level0.level1.level2.gross': {'$gt': 100000}}]}]}))

```



Experimental Results



Data diversity effects on query rewriting time and dictionary size

# of schemas	Query rewriting in (s)	Dictionary size
10	0.0005	40 KB
100	0.0025	74 KB
1 K	0.139	2 MB
3 K	0.6	7.2 MB
5 K	1.52	12 MB

Dictionary online construction overhead

#of schemas	Load (s)	Load and dict. (s)	Overhead
2	201s	269s	33%
4	205s	277s	35%
6	207s	285s	37%
8	208s	300s	44%
10	210s	309s	47%

Table: Study of the overhead added during load time

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Conclusion

EASYQ Advantages

- Overcoming the problem of querying documents with structural heterogeneity.
- Transparent rewriting mechanisms.
- Ensuring the coverage of latest structural changes. Therefore, the same query is rewritten at each execution

⇒ *The heterogeneity is automatically handled.*

Perspectives

- Employing real datasets
- Dealing with concurrent access
- Covering more operators

```

{''The_End'' :
  ''Thank you for your Kind Attention'' ,
  ''Next_?'' :
    ''It's Q&A Time '' ,
  ''Dataset'' :{
    ''Available_Online'':{
      }
    }
  }
}

```

